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$$\Pr[\text{LIN-CR}_\xi(G) \neq f(G)] + \Pr\left[|f(G) - \mathbb{E}[f(G)]| > 2\alpha e^{3/2}\right] \leq \exp(-e/4) + 2\exp(-\alpha^2/4).$$

If $\alpha \leq \sqrt{e}$, the last sum is at most $3\exp(-\alpha^2/4)$, as required. This concludes the proof of Theorem 7

Now we can prove Theorem 5. Fix $p = p(n)$ with $p(n) \gg \frac{\ln n}{n}$ and $G \sim G(n, p)$. Set $e = p\binom{n}{2}$. Let $C_n = \kappa_{\text{LIN-CR}}(n, 1)$. Since $\xi \in X$, $\mathbb{E}[\text{LIN-CR}_\xi(G)] = p^2 \text{LIN-CR}_\xi(K_n) \geq C_n e^2$. Let $\varepsilon > 0$ be arbitrarily small, but fixed. Then

$$\Pr\left[\text{LIN-CR}(G) < (C_n - \varepsilon)e^2\right] \leq \sum_{\xi \in X} \Pr\left[\text{LIN-CR}_\xi(G) < \mathbb{E}[\text{LIN-CR}_\xi(G)] - \varepsilon e^2\right]$$

We apply Theorem 7 with $3\alpha e^{3/2} = \varepsilon e^2$ so that $\alpha^2/4 = \frac{1}{36}\varepsilon^2 e$. The growth rate of $p(n)$ insures that this is $o(n^{6n})$ for any fixed positive ε . The Goodman-Pollack result critically bounds $|X| \leq n^{6n}$. Hence the sum goes to zero, as desired.

Comments and Open Questions: We have not been able to determine if the condition $p \gg \frac{\ln n}{n}$ in Theorem 5 is necessary. We have already conjectured that for any $p = p(n)$ with $np \rightarrow \infty$ we already have $\lim_{n \rightarrow \infty} \kappa_{\text{LIN-CR}}(n, p) = \gamma_{\text{LIN-CR}}$. While the Goodman-Pollack theorem itself cannot be improved asymptotically [2], it might be the case that there are few (in some sense) near optimal drawings so that the $n^{-\Theta(n)}$ error probability used in the proof of Theorem 5 may not be fully necessary. This, however, remains highly speculative.

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First, however, we examine a fixed order type $\xi \in X$. For any graph G with vertices v_1, \dots, v_n let G be a graph with vertices v_1, v_2, \dots, v_n . $\text{LIN-CR}_\xi(G)$ denote the number of crossings in the straight line drawing of G where v_i is placed at x_i in the plane and x_1, \dots, x_n have order type ξ .

Theorem 7 *Let $G(n, p)$ be a random graph with vertices v_1, v_2, \dots, v_n , with edge probability $0 < p = p(n) < 1$, and let $e = p\binom{n}{2}$. Then*

$$\Pr \left[|\text{LIN-CR}_\xi(G) - \mathbb{E}[\text{LIN-CR}_\xi(G)]| > 3\alpha e^{3/2} \right] < 3 \exp(-\alpha^2/4)$$

holds for every α satisfying $(e/4)^3 \exp(-e/4) \leq \alpha \leq \sqrt{e}$.

Proof: We follow the approach of Pach and Tóth [27]. (We note that general polynomial concentration results of Kim and Vu [19] could also be used.) Let $e_1, e_2, \dots, e_{\binom{n}{2}}$ be the edges of the complete graph on $V(G)$. Define another random graph G^* on the same vertex set, as follows. If G has at most $2e$ edges, let $G^* = G$. Otherwise, there is an $i < \binom{n}{2}$ so that $|\{e_1, e_2, \dots, e_i\} \cap E(G)| = 2e$, and set $E(G^*) = \{e_1, e_2, \dots, e_i\} \cap E(G)$. Finally, let $f(G) = \text{LIN-CR}_\xi(G^*)$.

The addition of any edge to G can modify the value of f by at most $2e$. Following the terminology of Alon–Kim–Spencer [3], we say that the *effect* of every edge is at most $2e$. The *variance* of any edge is defined as $p(1-p)$ times the square of its effect. Therefore, the *total variance* cannot exceed

$$\sigma^2 = \binom{n}{2} p(2e)^2 = 4e^3.$$

Applying the Martingale Inequality of [3], which is a variant of Azuma’s Inequality [5] (see also [4]), we obtain that for any positive $\alpha \leq \sigma/e = 2\sqrt{e}$,

$$\Pr \left[|f(G) - \mathbb{E}[f(G)]| > \alpha\sigma = 2\alpha e^{3/2} \right] < 2 \exp(-\alpha^2/4).$$

Our goal is to establish a similar bound for $\text{LIN-CR}_\xi(G)$ in place of $f(G)$. Obviously,

$$\Pr [f(G) \neq \text{LIN-CR}_\xi(G)] \leq \Pr [G \neq G^*] < \exp(-e/4).$$

Thus, we have

$$\begin{aligned} |\mathbb{E}[f(G)] - \mathbb{E}[\text{LIN-CR}_\xi(G)]| &\leq \Pr [f(G) \neq \text{LIN-CR}_\xi(G)] \max \text{LIN-CR}_\xi(G) \leq \\ &\exp(-e/4) \frac{n^4}{8} \leq \alpha e^{3/2}, \end{aligned}$$

whenever $\alpha \geq (e/4)^3 \exp(-e/4)$ (say). Therefore,

$$\Pr \left[|\text{LIN-CR}_\xi(G) - \mathbb{E}[\text{LIN-CR}_\xi(G)]| > 3\alpha e^{3/2} \right] \leq$$

$i^i 3^{ia/2} (\frac{A}{ia})^{ia/2}$. The probability of having these edges is at most $(c/n)^i (c/n)^{ia}$. Multiplying terms, the probability is less than

$$\left[\frac{ne}{i} \frac{c}{n} \left(\frac{c}{n} \right)^a i^i 3^{ia/2} \left(\frac{A}{ia} \right)^{a/2} \right]^i$$

which simplifies to

$$\left[ec^{1+a} A^{a/2} a^{-a/2} \left(\frac{i}{n} \right)^a \right]^i \leq \left[ec^{1+a} A^{a/2} a^{-a/2} \delta^a \right]^i$$

We select δ sufficiently small so that the bracketed term is less than one. Then the sum over $\delta n \geq i > \ln^{1/3} n$ is $o(1)$, completing the theorem.

Now we put the two results together to bound $\text{CR}(G)$. The classic Erdős-Rényi results give that $G(n, \frac{c}{n})$, with $c > 1$, almost surely have a giant component with $\sim kn$ vertices and $\sim kn(1+b)$ edges where k, b are explicit functions of c and both are positive. We split the vertex set into l parts so that the number of edges between vertices in different parts is $\leq (l-1)A(G)$. Each part has size $\leq \delta n$ where $\delta = \frac{8}{7}$. Pick l so large so that δ is so small that every $i \leq \delta n$ vertices have at most $i(1 + \frac{b}{2})$ edges. Then at most $kn(1 + \frac{b}{2})$ edges from the giant component are between vertices in the same part. Thus at least $\frac{kb}{2}n$ edges of the giant component are between vertices in different parts. That is, $(l-1)A(G) \geq \frac{kb}{2}n$. But l, k, b are all constants (i.e., dependent only on c) and so $A(G) = \Omega(n)$. Thus $\text{CR}(G) = \Omega(n^2)$.

Comments and Open Questions: From Theorem 1 we know $\kappa_{\text{CR}}(n, p) \rightarrow \gamma_{\text{CR}}$ as $p \rightarrow 1$ and we have just shown that $\kappa_{\text{CR}}(n, \frac{c}{n})$ is bounded from below. How large does $p = p(n)$ need to be so that $\kappa_{\text{CR}}(n, p(n)) \sim \gamma_{\text{CR}}$? We have already conjectured that for any $p = p(n)$ with $np \rightarrow \infty$ we have $\kappa_{\text{CR}}(n, p) \rightarrow \gamma_{\text{CR}}$. But we cannot even show that $\kappa_{\text{CR}}(n, p) \rightarrow \gamma_{\text{CR}}$ when $p < 1$ is a constant. Suppose (which is surely true though we are unable to show it) that $\lim_n \kappa_{\text{CR}}(n, \frac{c}{n})$ exists and call it $g(c)$. Then $g(c)$ would be increasing so $\lim_{c \rightarrow \infty} g(c)$ would exist but might be a value strictly less than γ_{CR} . Would there be a second (or even a third or more) region (something like $p = \Theta(n^{-1/2})$ or, more likely, $p = \Theta(1)$) where $\kappa_{\text{CR}}(n, p)$ increases (in some asymptotic sense) until it finally reaches γ_{CR} ?

5 The rectilinear crossing number

Here we show Theorem 5. An *order type* of the points x_1, x_2, \dots, x_n in the plane (with no three colinear) is a list of orientations of all triples $x_i x_j x_k$, $i < j < k$ [15]. Elementary geometry gives that the order type of the four triples $x_i x_j x_k, x_i x_j x_l, x_i x_k x_l, x_j x_k x_l$ determines whether or not the straight line segments $x_i x_j$ and $x_k x_l$ intersect. Let X be the set of all order types of the points x_1, x_2, \dots, x_n in the plane. We shall make critical use of a result of Goodman and Pollack [15, 16] that $|X| < n^{6n}$. We note that the Goodman-Pollack result is derived from the Milnor-Thom theorem, a now classical and very deep result concerning algebraic varieties.

is at most $A(G|_{V_i})$. Restricting a graph to a subset can only lower $\text{CR}(G)$ and lower the degrees so that $A(G|_{V_i}) \leq A(G)$. Replace V_i by T, B , setting $V_i \leftarrow T$, $V_{i+1} \leftarrow B$. Continue this procedure until the partition has l parts.

Observe that in the final partition $V = V_1 \cup \dots \cup V_l$ all the edges between V_i, V_j occurred exactly once as an edge from T to B . Hence the total number of edges between all V_i, V_j is at most $(l - 1)A(G)$.

What about the sizes of the V_i ? Here we use that at each stage we split the largest set. If a set V_i is split when the partition has i parts then $|V_i| \geq \frac{n}{i} \geq \frac{n}{l}$ and so each part has size at least $\frac{n}{3l}$. Any final set V_j must be created as a T or B at some time so it has size at least $\frac{n}{3l}$. For the upper bound (where 8 is surely not the best constant) we first show that for $l = 2^t$ all $|V_i| \leq \frac{4n}{l}$. This is immediate for $t = 0, 1, 2$. Suppose it is true for $t \geq 2$ with $l = 2^t$. At most $\frac{l}{2}$ sets have size greater than $\frac{2n}{l}$. Say $\frac{2n}{l} < |V_i| \leq \frac{4n}{l}$. After one split the smaller part has size $\leq \frac{2n}{l}$ and the larger has size $\leq \frac{8n}{3l}$. Splitting the larger set gives sets of size $\leq \frac{16n}{9l} \leq \frac{2n}{l}$. That is, it takes at most two splits to make all parts have size $\leq \frac{2n}{l}$. Thus in at most l splits this has been done (recall we always split the largest set) for all the at most $\frac{l}{2}$ such V_i . That is, by the time the partition has $2l$ sets every set has size at most $\frac{4n}{2l}$ completing the induction. For general l say $2^t \leq l < 2^{t+1}$. The split when there are l sets cannot have bigger sets than when there were 2^t sets so each has size $\leq \frac{4n}{2^t} \leq \frac{8n}{l}$.

Now we show a density result for small subgraphs of $G(n, \frac{c}{n})$ that uses a surprisingly subtle argument.

Theorem 6 *Fix $c > 0$ and $a > 0$. Then there exists $\delta > 0$ (dependent only on c, a so that $G(n, \frac{c}{n})$ almost surely has the following property: For every $0 < i \leq \delta n$ every set of i vertices has less than $i(1 + a)$ edges.*

Proof: If this property does not hold then there is an i with $0 < i \leq \delta n$ and a set of i vertices which form a *connected* subgraph with $\geq i(1 + a)$ edges. For i small we use the well known result that G almost surely has no bicyclic subgraphs. For completeness we give a very rough argument: There are $\binom{n}{i} \leq n^i$ choices for the i vertices, $\leq 2^{i(i-1)/2}$ choices (a gross overestimate) for a bicyclic subgraph on the i vertices and probability $(c/n)^{i+1}$ of having the i edges so the probability is bounded from above by $\frac{1}{n} c^{i+1} 2^{i(i-1)/2}$. Summing this for, say, $i \leq \ln^{1/3} n$ gives $o(1)$.

For $\delta n \geq i > \ln^{1/3} n$ we exert greater care. There are $\binom{n}{i} \leq (ne/i)^i$ choices for the i vertices. Now we look at the number of connected graphs on i vertices with $i(1 + a)$ edges. (Technically this is $\lfloor i(1 + a) \rfloor$ edges but as a is fixed and we've already taken care of the cases with i fixed this has negligible effect.) We use a result of Bollobás [9]: The number of connected graphs on m vertices with $m - 1 + j$ edges is at most $m^{m-2} m^{3j/2} (A/j)^{j/2}$. Here A is an absolute constant. (In an asymptotic sense Bollobás's work was greatly extended by Łuczak [23], who showed that $A = \frac{e}{12}$ is the best constant, and Bender, Canfield and McKay [6] who gave an asymptotic formula for this number valid when $m \rightarrow \infty$ through the entire range of j . For our purposes, however, it is more convenient to use Bollobás's result as it holds for *all* m, j .) In our case this factor is less than

$T \cup B$ such that $\frac{2}{3}|V| \geq |T|, |B| \geq \frac{1}{3}|V|$. (The specific constant $\frac{2}{3}$ is not essential here, we need only to assure that T, B are roughly the same size.) Leighton observed that there is an intimate relationship between the bisection width and the crossing number of a graph [21], which is based on the Lipton-Tarjan separator theorem for planar graphs [22]. The following version of this relationship was obtained by Pach, Shahrokhi, and Szegedy [25]. Let G be a graph on vertex set V with d_v denoting the degree of vertex v . Then

$$b(G) \leq 10\sqrt{\text{CR}(G)} + 2\sqrt{\sum_{v \in V(G)} d_v^2}$$

With $G \sim G(n, \frac{c}{n})$, $E[d_v^2] \sim c^2 + c = O(1)$ and almost surely $2\sqrt{\sum_{v \in V} d_v^2} = O(\sqrt{n})$ which proves to be negligible. For c a *large* constant basic probabilistic methods give that almost surely *every* partition $V = T \cup B$ with $\frac{2}{3}|V| \geq |T|, |B| \geq \frac{1}{3}|V|$ has a constant proportion of the edges running between them. That is, $b(G) = \Omega(n)$. Hence $\text{CR}(G) = \Omega(n^2)$.

Now suppose $c = 1 + \epsilon$ with $\epsilon > 0$ small. The difficulty is: almost surely $b(G)$ is zero! Why? From classic Erdős-Rényi results G will have a “giant component” of size $\sim kn$ with $k = k(c)$ and all other components will have size $O(\ln n)$. The function $k = k(c)$ was given explicitly by Erdős and Rényi but we need here only to note that $\lim_{c \rightarrow 1+} k(c) = 0$. For ϵ a small (actually, not so small) but fixed constant and $c = 1 + \epsilon$ the giant component has size kn with $k < \frac{2}{3}$. Place the giant component in the top T . Now take all other components sequentially. Add them to the top T if $|T|$ remains below $\frac{2}{3}n$, otherwise place them in the bottom B . This gives a partition with $\frac{2}{3}|V| \geq |T|, |B| \geq \frac{1}{3}|V|$ and no edges running between T and B .

With $c = 1 + \epsilon$, ϵ small, we shall split the vertices of G into a large but fixed number l of roughly equal parts. For convenience let us set

$$A(G) := 10\sqrt{\text{CR}(G)} + 2\sqrt{\sum_{v \in V(G)} d_v^2}$$

so that $b(G) \leq A(G)$. We claim that there exists a partition of the vertices V of G into l parts $V = V_1 \cup \dots \cup V_l$ such that (with $n = |V|$):

1. The total number of edges between vertices in different sets is at most $(l - 1)A(G)$
2. All $|V_i| \geq \frac{n}{3l}$
3. All $|V_i| \leq \frac{8n}{7}$

To achieve this partition we employ the following procedure. Begin with $V = V_1 \cup V_2$, $V_1 = T$, $V_2 = B$ as given by the Pach-Shahrokhi-Szegedy result [25]. Suppose at a general point we have a partition $V = V_1 \cup \dots \cup V_i$ where, renumbering for convenience, V_i is the largest set. Split $V_i = T \cup B$, using again the Pach-Shahrokhi-Szegedy result, so that the number of edges running from T to B

forming $K_5(L)$. Thus edges uv, wz lie on $\sim L^2 n^{10L-9} p^{10L-2}$ different $K_5(L)$. So each crossing has been counted at most that many times and hence the number of crossings is at least asymptotically

$$\frac{\frac{1}{5!} n^{10L-5} p^{10L}}{L^2 n^{10L-9} p^{10L-2}} = \frac{1}{120L^2} n^4 p^2$$

as desired.

Comments and Open Questions. As we must take $L > \varepsilon^{-1}$ the constant $\frac{1}{120}L^{-2}$ in this result goes to zero as $\varepsilon \rightarrow 0$. This is in surprising contrast to the crossing number $\text{CR}(G)$ discussed in the next section. That crossing number becomes a positive proportion of the square of the number of edges already at $p = \frac{c}{n}$ when $c > 1$. Can the pair-crossing number and the crossing number have such different behavior? We doubt it. As mentioned in the Introduction we cannot rule out the possibility that the pair-crossing number and the crossing number are always exactly the same. We can certainly make the weaker conjecture that the expectation of the pair-crossing number of $G(n, p)$ becomes $\Omega(n^4 p^2)$ already at $p = \frac{1+\varepsilon}{n}$. We further note that we have no idea at which p $\kappa_{\text{PAIR-CR}}(n, p)$ gets within $o(1)$ of its limit $\gamma_{\text{PAIR-CR}}$.

Pach and Tóth [28] introduced another variant of the crossing number. The *odd-crossing number* of any graph G is the minimum number of pairs of edges that cross an odd number of times, over all drawings of G . Clearly, for any graph, $\text{ODD-CR}(G) \leq \text{PAIR-CR}(G)$. With a little modification, the above argument works also for the odd-crossing number, therefore, the statement of Theorem 3 holds also for the odd-crossing number.

4 The crossing number

Here we prove Theorem 4. Fix $c > 1$ and set $p = \frac{c}{n}$. Let $G \sim G(n, p)$. Our object is to show

$$\liminf_{n \rightarrow \infty} \frac{E[\text{CR}(G)]}{\left[\binom{n}{2} p\right]^2} > 0$$

As c is constant this is equivalent to showing that for n sufficiently large

$$E[\text{CR}(G)] > \delta n^2$$

for some δ dependent only on c .

We begin by reviewing in outline form the argument of Pach and Tóth [28] which requires that c be a sufficiently large constant. We will see why their argument does not work for $c = 1 + \varepsilon$ with $\varepsilon > 0$ small and then how a modification of their argument, combined with results on $G(n, p)$, does work.

Define the *bisection width* of G , denoted by $b(G)$, as the minimal number of edges running between T (top) and B (bottom) over all partitions of the vertex set into two disjoint parts $V =$

This is equivalent to showing that for n sufficiently large

$$E[\kappa_{\text{PAIR-CR}}(G)] > \delta n^4 p^2$$

for some δ dependent only on ε . For $L \geq 1$ we let $K_5(L)$ denote the following graph:

- There are five vertices x_1, \dots, x_5
- For each distinct pair x_i, x_j there is a path between them of length L .

There are no other vertices nor edges so $K_5(L)$ has $5 + 10(L - 1)$ vertices and $10L$ edges. Note that $K_5(L)$ is a topological K_5 . Hence in any drawing of $K_5(L)$ there must be at least one crossing. We shall fix L such that $L\varepsilon > 1$. We shall show that G contains many $K_5(L)$. Each $K_5(L)$ will force at least one crossing. With L fixed this is a positive (albeit only $0.01L^{-2}$) proportion of the square of the number of edges involved. When this is carefully counted over all $K_5(L)$ we shall see that the total number of crossings is at least this constant times the square of the total number of edges.

We use three results about the almost sure behavior of $G(n, p)$. In the third K is any fixed constant.

1. Every vertex has degree $\sim np$.
2. Between every pair of distinct vertices there are $\sim n^{L-1}p^L$ paths of length L .
3. For any distinct x, y, z_1, \dots, z_K there are $\sim n^{L-1}p^L$ paths of length L between x and y that do not use any of the z_j .

The first result holds whenever $np \gg \ln n$ and follows from basic Large Deviation bounds on the degree of a vertex. Both the first and the second result are examples of a more general result [30] on counting extensions. For the third we note from [30] that the probability that the number of paths of length L between fixed x and y is not in $[(1 - \varepsilon)n^{L-1}p^L, (1 + \varepsilon)n^{L-1}p^L]$ is exponentially small. Fix x, y, z_1, \dots, z_K . Consider L -paths from x to y on G with z_1, \dots, z_K deleted, which has distribution $G(n - K, p)$. The K has negligible effect and so with exponentially small failure this number is as desired – hence almost surely it is as desired for all $O(n^{K+2})$ choices of x, y, z_1, \dots, z_K .

Now we count the $K_5(L)$. There are $\binom{n}{5} \sim \frac{1}{5!}n^5$ choices for the x_1, \dots, x_5 . Between each pair we place $\sim n^{L-1}p^L$ L -paths not using previously chosen paths. This gives a total of $\frac{1}{5!}n^{10L-5}p^{10L}$ copies of $K_5(L)$. For each one we count one crossing. Now consider a crossing between, say, edges uv and wz . How many $K_5(L)$ do they lie on? Renumbering for convenience say the path from x_1 to x_2 has u as its i -th and v as its $i + 1$ -st point and the path from x_3 to x_4 has w as its j -th and z as its $j + 1$ -st point. There are L^2 choices for i, j . Now fix u, v, w, z and i, j . From the first property there are $\sim (np)^i$ paths of length i starting at u , $\sim (np)^{L-i-1}$ paths of length $L - i$ starting at v and similarly for w, z . Further these numbers are not asymptotically effected when we require that they miss a fixed number of points. So we extend u, v, w, z to some x_1, x_2, x_3, x_4 in $\sim (np)^{2(L-1)}$ ways. We have n choices for x_5 and then $\sim (n^{L-1}p^L)^8$ ways to complete the remaining eight paths

particular crossing of the drawing of K_n in eight ways. so they have probability $8\text{LIN-CR}(K_n)/(n)_4$ of being mapped to a crossing. Now the expected number of crossings of G in this random drawing is at most, by Linearity of Expectation, $\frac{e^2 8\text{LIN-CR}(K_n)}{2(n)_4}$ and thus there exists a drawing of G with at most that many crossings.

As the right hand side approaches $\gamma_{\text{LIN-CR}}$ we have

$$\frac{\text{LIN-CR}(G)}{e^2} \leq \gamma_{\text{LIN-CR}} + o(1)$$

where the $o(1)$ term approaches zero in n , uniformly over all graphs G .

With $c > 0$ fixed (this argument is only needed for $c > 1$ but works for all positive c), $p = \frac{c}{n}$ and $G \sim G(n, p)$ let X denote the number of edges and Y denote the number of isolated edges. The savings comes from noting that isolated edges can always be added to a graph with no additional crossings. Thus

$$E[\text{LIN-CR}(G)] \leq E[(X - Y)^2](\gamma_{\text{LIN-CR}} + o(1))$$

Here $E[X] \sim \frac{c}{2}n$ and $E[Y] = \binom{n}{2}p(1-p)^{2n-4} \sim \frac{c}{2}e^{-2c}n$ and elementary calculations give

$$E[(X - Y)^2] \sim E[X - Y]^2 \sim \left[\frac{c}{2}(1 - e^{-2c})n\right]^2$$

With $e := p\binom{n}{2} \sim \frac{c}{2}n$ we have

$$\frac{E[\text{LIN-CR}(G)]}{e^2} \leq \gamma_{\text{LIN-CR}}(1 - e^{-2c})^2(1 + o(1))$$

Comments and Open Questions. We note that as c approaches infinity the $(1 - e^{-2c})^2$ term above approaches one. The above bound may be improved somewhat by letting Y denote the edges in isolated trees and unicyclic components and there are even further improvements possible. Still, all these improvements seem to approach one as c approaches infinity. This leads to an intriguing conjecture: If $p(n) \gg \frac{1}{n}$ then $\kappa_{\text{LIN-CR}}(n, p) \rightarrow \gamma_{\text{LIN-CR}}$. One may make the same conjecture for all three variants of the crossing number. Indeed, this entire paper may be viewed as an attempt (thus far unsuccessful) of the authors to resolve these conjectures.

We conjecture that for any $c \geq 0$, the limits $\lim_{n \rightarrow \infty} \kappa_{\text{LIN-CR}}(n, c/n)$, $\lim_{n \rightarrow \infty} \kappa_{\text{CR}}(n, c/n)$, and $\lim_{n \rightarrow \infty} \kappa_{\text{PAIR-CR}}(n, c/n)$ exist. This follows from Theorem 2, for $c \leq 1$. If this conjecture is true, it is not hard to see that the functions $f_{\text{LIN-CR}}(c) = \lim_{n \rightarrow \infty} \kappa_{\text{LIN-CR}}(n, c/n)$, $f_{\text{CR}}(c) = \lim_{n \rightarrow \infty} \kappa_{\text{CR}}(n, c/n)$, and $f_{\text{PAIR-CR}}(c) = \lim_{n \rightarrow \infty} \kappa_{\text{PAIR-CR}}(n, c/n)$ are continuous and increasing for all $c \geq 0$.

3 The pair-crossing number

Here we prove Theorem 3. Fix $\varepsilon > 0$ and set $p = p(n) = n^{\varepsilon-1}$. Our object is to show

$$\liminf_{n \rightarrow \infty} \kappa_{\text{PAIR-CR}}(n, p) > 0$$

with expected number of crossings $q^2 \kappa_{\text{CR}}(n, p)$. We do not claim this drawing is optimal, but it does give the desired upper bound as $E[\text{CR}(G(n, pq))] \leq q^2 E[\text{CR}(G(n, p))]$, completing Theorem 1.

The first six parts of Theorem 2 will come as no surprise to those familiar with random graphs as in the classic papers of Erdős and Rényi it was shown that with $p = \frac{c}{n}$ the random graph $G(n, p)$ is almost surely planar when $c < 1$. Our argument is a bit technical, however, as we must bound the expected crossing number.

Fix $c < 1$, set $p = \frac{c}{n}$ and $X = \text{LIN-CR}(G)$ with $G \sim G(n, p)$. Let Y be the number of cycles of G and Z the number of edges of G . Then we claim $X \leq YZ$. Remove from G one edge from each cycle. This leaves a forest which can be drawn with straight lines and no crossings. Now add back in those Y edges as straight lines. At worst they could hit every edge, giving $\leq YZ$ crossings. With $c < 1$ $E[Y] = \sum_{i=3}^n \frac{\binom{n}{i}}{2^i} p^i < \sum_{i=3}^{\infty} c^i$ is bounded by a constant, say A . As Z has Binomial Distribution standard bounds give, say, $\Pr[Z > 10n] < \alpha^{-n}$ for some explicit $\alpha > 1$. As $X \leq n^4$ always, $X \leq 10nY + n^4 \chi(Z > 10n)$ where χ is the indicator random variable. Thus $E[X] \leq 10An + n^4 \alpha^{-n} = o(n^2)$.

Now fix $c = 1 + \varepsilon$ with ε positive and small. Set $p = \frac{1+\varepsilon}{n}$, $p' = \frac{1-\varepsilon}{n}$ and let p^* satisfy $p' + p^* - p'p^* = p$ so that $p^* \sim \frac{2\varepsilon}{n}$. We may consider $G(n, p)$ as the union of independently chosen $G(n, p')$ and $G(n, p^*)$. Say the first has rectilinear crossing number X and Y edges and the second has Z edges. Then their union has rectilinear crossing number at most $X + Y(Y + Z)$ as we draw $G(n, p')$ optimally and assume all other pairs of edges do intersect. But $E[X] = o(n^2)$ and it is easy to show that $E(Y(Y + Z)) \sim E(Y)(E(Y + Z)) \sim \frac{1}{2}n^2\varepsilon(1 + \varepsilon)$. Thus

$$E[\text{LIN-CR}(G)] \leq (1 + o(1))\frac{1}{2}\varepsilon(1 + \varepsilon)n^2$$

from which part 4 of Theorem 2 follows. Parts 5 and 6 then also follow as they involve smaller crossing numbers.

The final three parts of Theorem 2 are also natural to those familiar with random graphs. For $c > 1$ fixed $G(n, \frac{c}{n})$ has a “giant component” with $\Omega(n)$ vertices. Outside the giant component there are $\Omega(n)$ edges all lying in trees or unicyclic components. These edges may be drawn with no crossings and that will involve a positive proportion of the potential edge crossings. Again, our argument will be a bit technical as we must deal with expectations. We state the argument only for rectilinear crossing number but it is the same in all three cases.

We first note a precise result: Let G be *any* graph on n vertices with e edges. Then

$$\frac{\text{LIN-CR}(G)}{e^2} \leq \frac{4\text{LIN-CR}(K_n)}{\binom{n}{4}}$$

Fix a drawing of K_n with $\text{LIN-CR}(K_n)$ crossings. Define a random drawing of G by randomly mapping its n vertices bijectively to the n vertices of the drawing. Let e_1, e_2 be two edges of G with no common vertex, there being at most $e^2/2$ such unordered pairs. They may be mapped to a

7. $\limsup_{n \rightarrow \infty} \kappa_{\text{LIN-CR}}(n, c/n) < \gamma_{\text{LIN-CR}}$ for all c
8. $\limsup_{n \rightarrow \infty} \kappa_{\text{CR}}(n, c/n) < \gamma_{\text{CR}}$ for all c
9. $\limsup_{n \rightarrow \infty} \kappa_{\text{PAIR-CR}}(n, c/n) < \gamma_{\text{PAIR-CR}}$ for all c

Theorem 2 gives only upper bounds for the various crossing numbers. The main results of this paper, given in Theorems 3, 4, 5, deal with lower bounds for the three crossing numbers. Our weakest result is for the pair-crossing number.

Theorem 3 For any $\varepsilon > 0$, $p = p(n) = n^{\varepsilon-1}$, $\liminf_{n \rightarrow \infty} \kappa_{\text{PAIR-CR}}(n, p) > 0$.

For the crossing number we have a much stronger result.

Theorem 4 For any $c > 1$ with $p = p(n) = c/n$

$$\liminf_{n \rightarrow \infty} \kappa_{\text{CR}}(n, p) > 0$$

As $\text{LIN-CR}(G) \geq \text{CR}(G)$ the lower bound of Theorem 4 applies also to the rectilinear crossing number. Our most surprising result is that with the rectilinear crossing number one reaches an asymptotically best limit in relatively short time.

Theorem 5 If $p = p(n) \gg \frac{\ln n}{n}$ then

$$\lim_{n \rightarrow \infty} \kappa_{\text{LIN-CR}}(n, p) = \gamma_{\text{LIN-CR}}(n, p)$$

2 Upper Bounds

Let f be any real valued function on graphs. Then with $G \sim G(n, p)$

$$E[f(G)] = \sum_H f(H) p^{e(H)} (1-p)^{\binom{n}{2} - e(H)}$$

where H runs over the labelled graphs on n vertices and $e(H)$ is the number of edges of H . This is a polynomial and hence a continuous function of p , giving the first part of Theorem 1. We argue that $\kappa_{\text{CR}}(n, p)$ is an increasing function of p , the other arguments being identical. For $0 \leq p, q \leq 1$ we may view $G(n, pq)$ as a two step process, first creating $G(n, p)$ and then taking each edge from $G(n, p)$ with probability q . After the first stage consider a drawing with the minimal number of crossing X , so that $E[X] = \kappa_{\text{CR}}(n, p)$. Now keep that drawing but take each edge with probability q . Each crossing is still in the new picture with probability q^2 . This gives a drawing of $G(n, pq)$

These limits are known to exist [29] and the best known bounds are $1/30 \leq \gamma_{\text{PAIR-CR}} \leq 1/16$, $1/20 \leq \gamma_{\text{CR}} \leq 1/16$, $1/20 \leq \gamma_{\text{LIN-CR}} \leq 0.639$ [18, 29, 10].

In this paper we investigate the crossing numbers of *random graphs*. Let $G = G(n, p)$ be a *random graph* with n vertices, whose edges are chosen independently with probability p . Let e denote the *expected number* of edges of G , i.e., $e = p\binom{n}{2}$. We shall always have $e \rightarrow \infty$ (indeed, $p = \Omega(n^{-1})$) so that G almost surely has $e(1 + o(1))$ edges.

In [27] it was shown that if $e > 10n$, then almost surely we have $\text{CR}(G) \geq \frac{e^2}{4000}$. Consequently, almost surely we also have $\text{LIN-CR}(G) \geq \frac{e^2}{4000}$. As we always *can* draw with straight lines the crossing number (in any form) is never larger than the number of pairs of edges and the expected number of pairs of edges is $\frac{e^2}{2}$. Our interest will be in those regions of p for which the various crossing numbers are, asymptotically, a positive proportion of the number of pairs of edges.

Let

$$\kappa_{\text{LIN-CR}}(n, p) = \frac{\mathbb{E}[\text{LIN-CR}(G)]}{e^2}, \quad \kappa_{\text{CR}}(n, p) = \frac{\mathbb{E}[\text{CR}(G)]}{e^2}, \quad \kappa_{\text{PAIR-CR}}(n, p) = \frac{\mathbb{E}[\text{PAIR-CR}(G)]}{e^2}.$$

We have $\kappa_{\text{PAIR-CR}}(n, p) \leq \kappa_{\text{CR}}(n, p) \leq \kappa_{\text{LIN-CR}}(n, p)$ for any n, p .

Theorem 1 *For any fixed n , $\kappa_{\text{LIN-CR}}(n, p)$, $\kappa_{\text{CR}}(n, p)$, $\kappa_{\text{PAIR-CR}}(n, p)$ are increasing, continuous functions of p .*

With Theorem 1 we may express (roughly) our two central concerns. At which $p = p(n)$ are $\kappa_{\text{LIN-CR}}(n, p)$, $\kappa_{\text{CR}}(n, p)$, $\kappa_{\text{PAIR-CR}}(n, p)$ bounded away from zero? At which $p = p(n)$ are $\kappa_{\text{LIN-CR}}(n, p)$, $\kappa_{\text{CR}}(n, p)$, $\kappa_{\text{PAIR-CR}}(n, p)$ close to their limiting values $\gamma_{\text{LIN-CR}}$, γ_{CR} , $\gamma_{\text{PAIR-CR}}$? Our results for these three variants of crossing number shall be quite different. We are uncertain whether or not that represents the reality of the situation. The following relatively simple result shows basically that for $p = \frac{1}{n}$ all three crossing numbers are asymptotically negligible and that for $p = \frac{c}{n}$ with $c > 1$ fixed the three crossing numbers have not reached their limiting values.

Theorem 2 1. $\limsup_{n \rightarrow \infty} \kappa_{\text{LIN-CR}}(n, c/n) = 0$ for $c \leq 1$

2. $\limsup_{n \rightarrow \infty} \kappa_{\text{CR}}(n, c/n) = 0$ for $c \leq 1$

3. $\limsup_{n \rightarrow \infty} \kappa_{\text{PAIR-CR}}(n, c/n) = 0$ for $c \leq 1$

4. $\lim_{c \rightarrow 1} \limsup_{n \rightarrow \infty} \kappa_{\text{LIN-CR}}(n, c/n) = 0$

5. $\lim_{c \rightarrow 1} \limsup_{n \rightarrow \infty} \kappa_{\text{CR}}(n, c/n) = 0$

6. $\lim_{c \rightarrow 1} \limsup_{n \rightarrow \infty} \kappa_{\text{PAIR-CR}}(n, c/n) = 0$

Crossing numbers of random graphs

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Abstract

The *crossing number* of G is the minimum number of crossing points in any drawing of G . We consider the following two other parameters. The *rectilinear crossing number* is the minimum number of crossing points in any drawing of G , with straight line segments, as edges. The *pairwise crossing number* of G is the minimum number of pairs of edges that cross over all drawings of G . We prove several results on the expected values of these parameters of a random graph.

1 Introduction

A *drawing* of a graph G is a mapping which assigns to each vertex a point of the plane and to each edge a simple continuous arc connecting the corresponding two points. We assume that in a drawing no three edges (arcs) cross at the same point, and the edges do not pass through any vertex. The *crossing number* $\text{CR}(G)$ of G is the minimum number of crossing points in any drawing of G . We consider the following two variants of the crossing number. The *rectilinear crossing number* is the minimum number of crossing points in any drawing of G , with straight line segments, as edges. The *pairwise crossing number* $\text{PAIR-CR}(G)$ of G is the minimum number of pairs of edges that cross over all drawings of G .

Clearly, $\text{PAIR-CR}(G) \leq \text{CR}(G) \leq \text{LIN-CR}(G)$.

Bienstock and Dean [8] constructed a series of graphs with crossing number 4, whose rectilinear crossing numbers are arbitrary large. On the other hand, Pach and Tóth [28] proved that for any graph G , $\text{CR}(G) \leq 2(\text{PAIR-CR}(G))^2$. Probably this bound is very far from being optimal, we can not even rule out that $\text{CR}(G) = \text{PAIR-CR}(G)$ for any graph G .

The determination of the crossing numbers is extremely difficult. Even the crossing numbers of the complete graphs are not known. Let

$$\gamma_{\text{PAIR-CR}} = \lim_{n \rightarrow \infty} \frac{\text{PAIR-CR}(K_n)}{\binom{n}{2}^2}, \quad \gamma_{\text{CR}} = \lim_{n \rightarrow \infty} \frac{\text{CR}(K_n)}{\binom{n}{2}^2}, \quad \gamma_{\text{LIN-CR}} = \lim_{n \rightarrow \infty} \frac{\text{LIN-CR}(K_n)}{\binom{n}{2}^2}.$$

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