Ordered graphs and large bi-cliques in intersection graphs of curves

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Abstract

An *ordered graph* $G_<$ is a graph with a total ordering $<$ on its vertex set. A *monotone path* of length $k-1$ is a sequence of vertices $v_1 < v_2 < \ldots < v_k$ such that v_iv_j is an edge of $G_<$ if and only if $|j-i| = 1$. A *bi-clique* of size m is a complete bipartite graph whose vertex classes are of size m.

We prove that for every positive integer k, there exists a constant $c_k > 0$ such that every ordered graph on n vertices that does not contain a monotone path of length k as an induced subgraph has a vertex of degree at least $c_k n$, or its complement has a bi-clique of size at least $c_k n/\log n$. A similar result holds for ordered graphs containing no induced ordered subgraph isomorphic to a fixed ordered matching.

As a consequence, we give a short combinatorial proof of the following theorem of Fox and Pach. There exists a constant $c > 0$ such the intersection graph G of any collection of n x-monotone curves in the plane has a bi-clique of size at least $cn/\log n$ or its complement contains a bi-clique of size at least cn. (A curve is called x-monotone if every vertical line intersects it in at most one point.) We also prove that if G has at most $\left(\frac{1}{4}-\epsilon\right)\binom{n}{2}$ edges for some $\epsilon > 0$, then $\overline{G}$ contains a linear sized bi-clique. We show that this statement does not remain true if we replace $\frac{1}{4}$ by any larger constants.

1 Introduction

There are a growing number of examples showing that ordered structures can be useful for solving geometric and topological problems that appear to be hard to analyze by traditional combinatorial methods. The aim of the present paper is to provide an example concerning intersection patterns of curves, where one can apply ordered graphs.

First, we agree on the terminology. An *ordered graph* $G_<$ is a graph G with a total ordering $<$ on its vertex set. If the ordering $<$ is clear from the context, we write G instead of $G_<$. An ordered graph $H_{<'}$ is an *induced subgraph* of the ordered graph $G_<$, if there exists an embedding $\phi : V(H) \to V(G)$ such that for every $u, v \in V(H)$, if $u <' v$ then $\phi(u) < \phi(v)$, and $uv \in E(H)$ if and only if $\phi(u)\phi(v) \in E(G)$.

A *monotone path* P_k of length $k-1$ is an ordered graph with k vertices $v_1 < v_2 < \ldots < v_k$ in which v_iv_j is an edge if and only if $|j-i| = 1$. A *bi-clique* in an (ordered or unordered) graph G consists of a pair of disjoint subsets of the vertices (A, B) such that $|A| = |B|$ and for every $a \in A$ and $b \in B$, there is an edge between a and b. The size of a bi-clique (A, B) is $|A|$. A *comparability*

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graph is a graph G for which there exists a partial ordering on $V(G)$ such that two vertices are joined by an edge of G if and only if they are comparable by this partial ordering. An *incomparability graph* is the complement of a comparability graph. The maximum degree of the vertices of G is denoted by $\Delta(G)$.

Our first theorem states that if a P_k-free ordered graph is not too dense, then its complement contains a large bi-clique.

Theorem 1. *For every integer $k \geq 2$, there exists a constant $c = c(k) > 0$ such that the following statement is true. Let $G_<$ be an ordered graph on n vertices which satisfies $\Delta(G_<) < cn$ and does not have any induced ordered subgraph isomorphic to the monotone path P_k.*

Then the complement of $G_<$ contains a bi-clique of size at least $cn/\log n$.

For the conclusion to hold, we need some upper bound on the degrees of the vertices (or on the number of edges) of the graph. To see this, consider the graph G on the naturally ordered vertex set $\{1, \ldots, n\}$, in which $A = \{1, \ldots, \lfloor n/2 \rfloor\}$ and $B = \{\lfloor n/2 \rfloor + 1, \ldots, n\}$ induce complete subgraphs, and any pair of vertices $a \in A, b \in B$ are joined by an edge randomly, independently with a very small probability $p > 0$. This ordered graph has no induced monotone path of length 4, its maximum degree satisfies $\Delta(G) < (1/2 + p)n$, but the maximum size of a bi-clique in its complement is $O_p(\log n)$. Consequently, for the constant appearing in Theorem 1, we have $c_5 \leq 1/2$.

The assumption that $G_<$ contains no induced P_3 is equivalent to the property that $G_<$ is a comparability graph. In this special case (that is, for $k = 2$), Theorem 1 was established by Fox, Pach, and Tóth [9], and in a weaker form by Fox [6]. Apart from the value of the constant c, the bound is best possible for $k = 2$ and, hence, for every $k \geq 2$.

An *ordered matching* is an ordered graph on $2k$ vertices which consists of k edges, no two of which share an endpoint. Our next result is an analogue of Theorem 1 for ordered graphs that contain no induced subgraph isomorphic to a fixed ordered matching.

Theorem 2. *For every ordered matching M, there exists a constant $c = c(M) > 0$ such the following statement is true. Let $G_<$ be an ordered graph on n vertices which satisfies $\Delta(G_<) < cn$ and does not have any induced ordered subgraph isomorphic to M.*

Then the complement of $G_<$ contains a bi-clique of size at least cn.

The conclusion of Theorem 2 is stronger than that of Theorem 1: in this case we can find a linear-sized bi-clique in the complement of $G_<$.

Given a family of sets, $\mathcal{C}$, the *intersection graph* of $\mathcal{C}$ is the graph, whose vertices correspond to the elements of $\mathcal{C}$, and two vertices are joined by an edge if and only if the corresponding sets have a nonempty intersection. A *curve* is the image of a continuous function $\phi : [0, 1] \to \mathbb{R}^2$. A curve is said to be *x-monotone* if every vertical line intersects it in at most one point. Note that any convex set can be approximated arbitrarily closely by x-monotone curves, so the notion of x-monotone curve extends the notion of convex sets. Throughout this paper, a curve will be called a *grounded* if one of its endpoints lies on the y-axis (on the vertical line $\{x = 0\}$) and the whole curve is contained in the nonnegative half-plane $\{x \geq 0\}$. (By slight abuse of notation, we write $\{x \geq 0\}$ for the set $\{(x, y) \in \mathbb{R}^2 : x \geq 0\}$.)

We will apply Theorems 1 and 2 to give a simple combinatorial proof for the following Ramsey-type result of Fox, Pach and Tóth [9], which is related to a celebrated conjecture of Erdős and Hajnal [4, 2].

Theorem 3. *[9] There exists an absolute constant $c > 0$ with the following property. The intersection graph G of any collection of n x-monotone curves contains a bi-clique of size at least $cn/\log n$, or its complement $\overline{G}$ contains a bi-clique of size at least cn.*

This result is tight, up to the value of c; see [16]. Indeed, Fox [6] proved that for any $\varepsilon > 0$ there exists a constant $c(\varepsilon)$ such that for every $n \in \mathbb{N}$, there exists an incomparability graph G on n vertices such that G does not contain a bi-clique of size $c(\varepsilon)n/\log n$, and the complement of G does not contain a bi-clique of size n^{ϵ}. On the other hand, every incomparability graph is isomorphic to the intersection graph of a collection of x-monotone curves [20, 11, 16].

It was shown in [9] that if the intersection graph of n x-monotone curves has at most $12^{-8}\binom{n}{2}$ edges, then the second option holds in Theorem 3: $\overline{G}$ contains a bi-clique of size at least cn. Also, the same result (with different constants) follows from a separator theorem of Lee [10] for string graphs. None of these arguments leave much room for replacing 12^{-8} by a decent constant. Tomon [21] applied some properties of partially ordered sets to establish the upper bound $\left(\frac{1}{16} - o(1)\right)\binom{n}{2}$. Somewhat surprisingly, using ordered graphs, one can precisely determine the best constant for which the statement still holds.

Theorem 4. *For any $\epsilon > 0$, there are constants $c_1 = c_1(\epsilon), c_2 = c_2(\epsilon) > 0$, and an integer $n_0 = n_0(\epsilon)$ such that the following statements are true. For every $n \geq n_0$,*

(1) there exist n x-monotone curves such that their intersection graph G has at most $(\frac{1}{4} + \epsilon)\binom{n}{2}$ edges, but the complement of G does not contain a bi-clique of size $c_1 \log n$;

(2) for any n x-monotone curves such that their intersection graph G has at most $(\frac{1}{4} - \epsilon)\binom{n}{2}$ edges, the complement of G contains a bi-clique of size $c_2 n$.

It is easy to see that every intersection graph of *convex sets* in the plane is also an intersection graph of x-monotone curves. We prove (1) by constructing n convex sets in the plane whose intersection graphs meets the requirements. Therefore, $\frac{1}{4}\binom{n}{2}$ is also a threshold for the emergence of linear sized bi-cliques in the complements of intersection graphs of convex sets.

It follows from the combination of Corollary 1.2 in [8] and the separator theorem of Lee [10] that Theorem 3 is true in a more general setting: without assuming that the curves are x-monotone. It is a serious challenge to extend our proof to that case. We still believe that Theorem 4 should also generalize to arbitrary curves.

Conjecture 5. *For any $\epsilon > 0$, there exist $c_0 = c_0(\epsilon) > 0$ and $n_0 = n_0(\epsilon)$ with the property that for any collection of $n \geq n_0$ curves whose intersection graph has at most $(\frac{1}{4} - \epsilon)\binom{n}{2}$ edges, the complement of G contains a bi-clique of size $c_0 n$.*

For *unordered* graphs without (unordered) induced paths of length $k - 1$, the size of the largest bi-clique that can be found in $\overline{G}$ is larger than what was shown in Theorem 1: it is linear in n. More precisely, Bousquet, Lagoutte, and Thomassé [1] proved that for every positive integer k, there exists $c(k) > 0$ such that, if G is an unordered graph with n vertices and at most $c(k)\binom{n}{2}$ edges, which does

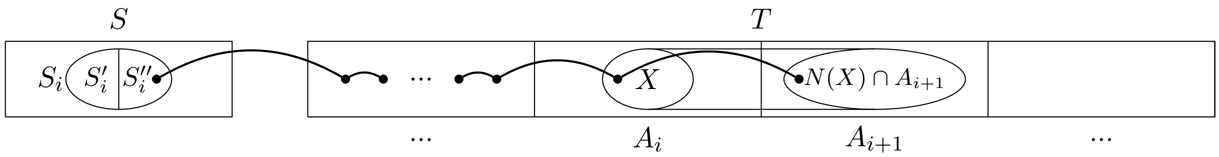

Figure 1: An illustration for the proof of Lemma 6.

not have an induced path of length $k-1$, then its complement $\overline{G}$ contains a bi-clique of size at least $c(k)n$. Recently, Chudnovsky, Scott, Seymour, and Spirkl [3] generalized this result to any forbidden forest, instead of a path. In an upcoming work [14], we obtain similar extensions of Theorems 1 and 2 to other *ordered forests*.

Our paper is organized as follows. Theorems 1 and 2 are proved in Sections 2 and 3, respectively. In Section 4, we first establish Theorem 3 for *grounded* x-monotone curves, and then show that this already implies the general result. Finally, we prove Theorem 4 in Section 5.

2 Ordered graphs avoiding a monotone induced path –Proof of Theorem 1

For any subset U of the vertex set of a graph G, define the *neighborhood* of U, as

$$N(U) = \{v \in V(G) \setminus U : \exists u \in U \text{ such that } uv \in E(G)\}.$$

If U consists of a single point u, we write $N(u)$ instead of $N(\{u\})$. The subgraph of G *induced* by the vertices in U is denoted by $G[U]$.

Given an ordered graph $G = G_<$ and two subsets $S, T \subset V(G)$, we write $S < T$ if $s < t$ for every $s \in S$ and $t \in T$. We say that a vertex $t \in T$ can be *reached* from a vertex $v \in V(G)$ *by a monotone T-path,* if there is an increasing sequence of vertices $v < t_1 < t_2 < \ldots < t_r = t$ such that $t_1, \ldots, t_r \in T$ and $vt_1, t_1t_2, \ldots, t_{r-1}t_r \in E(G)$. (The vertex v does not necessarily belong to T.) Let $P_G(S,T)$ denote the set of vertices in T that can be reached from *some* vertex in S by a monotone T-path in G. If it is clear from the context what the underlying ordered graph G is, we write $P(S,T)$ instead of $P_G(S,T)$. If S consists of a single vertex s, we write $P(s,T)$ instead of $P(\{s\},T)$. Finally, if $T = V(G)$, then we write $P(S)$ instead of $P(S, V(G))$.

For the proof of Theorem 1, we need the following lemma.

Lemma 6. *Let $G = G_<$ be an ordered graph on the vertex set $S \cup T$, where $S < T$, $|S| \geq \frac{n}{6\log_2 n}$, and $|T| \geq n$. Then either there exists a vertex $v \in S$ such that $|P(v,T)| \geq \frac{n}{12}$ or the complement of G contains a bi-clique of size $\frac{n}{12\log_2 n}$.*

Proof. With no danger of confusion, we omit the use of floors and ceilings, whenever they are not crucial. Let $m = 2^k$ such that $\frac{n}{12\log_2 n} < m \leq \frac{n}{6\log_2 n}$, and suppose that $\overline{G}$, the complement of G, contains no bi-clique of size m. Divide T into $s = \frac{n}{3m} \geq 2\log_2 n$ intervals of size $3m$, denoted by $A_1, \ldots, A_s$, in this order.

We will recursively define a sequence of sets $S \supset S_0 \supset S_1 \supset \cdots \supset S_k$ such that $|S_i| = 2^{k-i}$ for $i = 0, \ldots, k$, and $|P(S_i, T) \cap A_i| \geq m$ for $i = 1, \ldots, k$.

Let S_0 be an arbitrary m element subset of S. Suppose that the set S_i satisfying the above conditions has already been determined for some $i < k$. We define S_{i+1}, as follows. Let $X = P(S_i, T) \cap A_i$ if $i \geq 1$, and $X = S_0$ if $i = 0$. We can assume that $|N(X) \cap A_{i+1}| \geq 2m$, otherwise $|A_{i+1} \setminus N(X)| \geq m$ and there is a bi-clique (A, B) in $\overline{G}$ of size m such that $A \subset X$ and $B \subset A_{i+1} \setminus N(X)$. As A_{i+1} is $<$-larger than every element of A_i, we have $|N(X) \cap A_{i+1}| \subset P(S_i, T) \cap A_{i+1}$ and, hence, $|P(S_i, T) \cap A_{i+1}| \geq 2m$.

Partition S_i arbitrarily into two sets of size $|S_i|/2$, denoted by S' and S''. Clearly, we have $P(S', T) \cup P(S'', T) = P(S_i, T)$, so either $|P(S', T) \cap A_{i+1}| \geq m$, or $|P(S'', T) \cap A_{i+1}| \geq m$. In the first case, set $S_{i+1} = S'$, in the second, set $S_{i+1} = S''$. See Figure 1 for an illustration.

At the end of the process, S_k consists of one vertex, say v, and $|P(v, T) \cap A_k| \geq m$. We can actually assume that $|P(v, T) \cap A_j| \geq m$ for $j = k + 1, \ldots, s$. Indeed, there is no edge between $P(v, T) \cap A_k$ and $A_j \setminus P(v, T)$, so if $|A_j \cap P(v, T)| \leq m$, then $|A_j \setminus P(v, T)| \geq 2m$, which means that there exists a bi-clique (A, B) of size m in $\overline{G}$ such that $A \subset P(v, T) \cap A_k$ and $B \subset A_j \setminus P(v, T)$.

Summing up, we obtain that

$$|P(v, T)| \geq \sum_{j=k+1}^{s} |P(v, T) \cap A_j| \geq m(s - k) > \frac{n}{12}.$$

□

Proof of Theorem 1. Let $c = 1/(24k^2)$. Let $G = G_<$ be an ordered graph on n vertices of maximum degree at most cn, and suppose that $\overline{G}$ does not contain a bi-clique of size $m = cn/\log_2 n$. We have to prove that $G_<$ contains P_k as an induced subgraph.

Let the vertex set of G be $\{1, \ldots, n\}$, and for $i = 1, \ldots, k$, let $A_i = \{(i-1)n/k + 1, \ldots, in/k\}$. We recursively construct a sequence of vertices $x_1 < \cdots < x_k$ that satisfy conditions (1) and (2) below, for $l = 1, \ldots, k$. Let

$$U_{l+1} = V(G) \setminus \left(\bigcup_{i=1}^{l-1} N(x_i) \right).$$

Then

(1) $\{x_1, \ldots, x_l\}$ is an induced copy of P_l,

(2) $|P(x_l, U_{l+1}) \cap A_{l+1}| \geq m$.

For $l = 1$, apply Lemma 6 to the subgraph of G induced by $A_1 \cup A_2$ with $S = A_1$, $T = A_2$, and n/k instead of n. Then there exists $x_1 \in A_1$ such that $|P_G(x_1) \cap A_2| \geq \frac{n}{12k} > m$.

Now let $l > 1$ and suppose that the vertices $x_1 < \cdots < x_{l-1}$ satisfying conditions (1) and (2) have already been defined. Let $S = P(x_{l-1}, U_l) \cap A_l$ and $T = U_{l+1} \cap A_{l+1}$. (Note that for the definition of U_{l+1} we do not need x_l.) Then $|S| \geq m$ and, as the maximum degree of G is at most cn, we have $|T| \geq |A_{l+1}| - (l - 1)cn > \frac{n}{2k}$. Apply Lemma 6 to the subgraph of G induced by $S \cup T$ with $n/2k$ instead of n. Since $\overline{G}$ does not contain a bi-clique of size

$$\frac{cn}{\log_2 n} < \frac{n/(2k)}{12 \log_2(n/(2k))},$$

there exists $w \in S$ such that $|P(w,T)| > \frac{n}{24k}$. We have $w \in P(x_{l-1}, U_l)$, therefore w can be reached from x_{l-1} by a monotone U_l-path. Let $x_{l-1} = u_0 < \cdots < u_r = w$ be such a path with the minimum number of vertices. By the definition of U_l, the vertices $u_1, \ldots, u_r \in U_l$ do not belong to the neighborhoods of $x_1, \ldots, x_{l-2}$, and, by the minimality of the path, $u_2, \ldots, u_r$ are not in the neighborhood of x_{l-1}. Setting $x_l = u_1$, we find that $\{x_1, \ldots, x_l\}$ is an induced copy of P_l. Thus, condition (1) is satisfied.

The vertex w can be reached from x_l by the monotone U_{l+1}-path $x_l = u_1, \ldots, u_r = w$, and every $z \in P(w,T)$ can be reached from w by a monotone U_{l+1}-path. Therefore, every $z \in P(w,T)$ can be reached from x_l by a monotone U_{l+1}-path. This yields that

$$|P(x_l, U_{l+1}) \cap A_{l+1}| \geq |P(w,T)| > \frac{n}{24k} > m,$$

so that condition (2) is satisfied.

For $l = k$, the ordered subgraph of G induced by $\{x_1, \ldots, x_k\}$ is isomorphic to P_k. This completes the proof of the theorem. □

3 Ordered graphs avoiding an induced matching –Proof of Theorem 2

Proof of Theorem 2. Let k be the number of edges of M, and set $c = 1/(8k^2)$. Let $G = G_<$ be an ordered graph on n vertices such that the maximum degree of G is at most cn, and suppose that $\overline{G}$ does not contain a bi-clique of size cn. We have to prove that G contains M as an induced subgraph.

Suppose that $\{1, \ldots, 2k\}$ is the vertex set of M, and let $\{a_1, b_1\}, \ldots \{a_k, b_k\}$ be the edges of M. Let the vertex set of G be $\{1, \ldots, n\}$, and let $A_1, \ldots, A_{2k}$ be a partition of $V(G)$ into $2k$ intervals of size $\frac{n}{2k}$. Observe that, for every $i = 1, \ldots, k$, there exists a matching E_i of size at least $\frac{n}{4k}$ between A_{a_i} and A_{b_i}. Indeed, if F is a maximal sized matching between A_{a_i} and A_{b_i}, then $(A_{a_i} \setminus V(F), A_{b_i} \setminus V(F))$ is a bi-clique in $\overline{G}$ of size $\frac{n}{2k} - |F|$, so $|F| > \frac{n}{4k}$.

We show that for $i = 1, \ldots, k$, we can pick an edge $e_i = \{u_{a_i}, u_{b_i}\} \in E_i$ such that $\{u_j, u_l\}$ is not an edge if $\{j, l\} \notin E(M)$. If this is true, then $\{u_1, \ldots, u_{2k}\}$ spans a copy of M in G, and we are done. Pick the edges $e_1, \ldots, e_k$ one-by-one, that is, if $e_1, \ldots, e_i$ are already defined for some $i < k$ such that $\{e_1, \ldots, e_i\}$ is an induced matching in G, then define e_{i+1} as follows. The number of vertices in G that are in the neighborhood of any endpoint of the edges $e_1, \ldots, e_i$ is at most $2i\Delta(G) < 2kcn = |E_{i+1}|$. But then there exists an edge $e_{i+1} \subset E_{i+1}$ such that $\{e_1, \ldots, e_{i+1}\}$ is also an induced matching in G. □

4 Intersection graphs of curves–Proof of Theorem 3

First, we prove Theorem 3 in the special case where the curves are grounded, that is, their left endpoints lie on the y-axis.

Lemma 7. *There exists an absolute constant $c > 0$ with the following property. The intersection graph G of any collection $\mathcal{C}$ of n grounded x-monotone curves contains a bi-clique of size at least $cn/\log n$, or its complement $\overline{G}$ contains a bi-clique of size at least cn.*

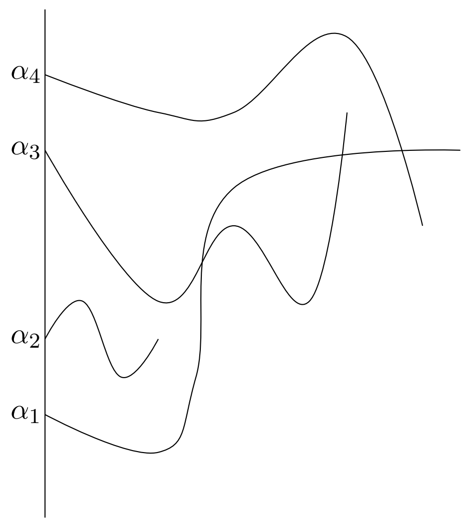

Figure 2: An illustration for the proof of part (2) of Lemma 8.

To prove this lemma, we first show that the intersection graphs of any collection of grounded x-monotone curves can be ordered in such a way that it has no induced ordered matching consisting of two crossing edges, or its complement has no monotone path P_4.

Let M_1 denote the ordered matching on vertex set $\{1, 2, 3, 4\}$, with edges $\{1, 3\}$ and $\{2, 4\}$.

Lemma 8. *Let $\mathcal{C}$ be a family of grounded curves (not necessarily x-monotone), let G be the intersection graph of $\mathcal{C}$, and let $<$ be the total ordering of $\mathcal{C}$ according to the y-coordinates of the endpoints of the elements of $\mathcal{C}$ lying on $\{x = 0\}$.*

(1) $G_<$ does not contain M_1 as an induced subgraph.

(2) If, in addition, the elements of $\mathcal{C}$ are x-monotone curves, then $\overline{G}_<$ does not contain P_4 as an induced subgraph.

Proof. (1) Suppose that $G_<$ contains M_1 as an induced subgraph, and let $\alpha_1 < \alpha_2 < \alpha_3 < \alpha_4$ denote the curves corresponding to the vertices of M_1. As α_1 and α_3 intersect, the line $\{x = 0\}$ and the two curves α_1, α_3 enclose a closed bounded region A. Curve α_2 is disjoint from both α_1 and α_3, and its endpoint on $\{x = 0\}$ belongs to A, we have $\alpha_2 \subset A$. Curve α_4 is also disjoint from α_1 and α_3, but its endpoint on $\{x = 0\}$ is not in A, so $\alpha_4 \cap A = \emptyset$. Hence, α_2 and α_4 cannot intersect, contradiction.

(2) Suppose that $\overline{G}_<$ contains P_4 as an induced subgraph, and let $\alpha_1 < \alpha_2 < \alpha_3 < \alpha_4$ denote the corresponding vertices. As before, the line $\{x = 0\}$ and the two curves α_1, α_3 enclose a closed bounded region A, and $\alpha_2 \subset A$. Since α_1 and α_3 are x-monotone, every vertical line intersecting A intersects α_1 and α_3 in exactly one point, the intersection point with α_1 lying below the intersection point with α_3. Curve α_4 is disjoint from α_3, so for every vertical line intersecting α_3 and α_4, its intersection with α_3 is below its intersection with α_4. Therefore, we have $A \cap \alpha_4 = \emptyset$, which implies that α_2 and α_4 are disjoint, contradiction. See Figure 2 for an illustration. $\square$

In view of Lemma 8, we may be able to use Theorem 1 or Theorem 2 to argue that G, the intersection graph of a collection of n x-monotone curves, contains a bi-clique of size $\Omega(n/\log n)$, or its complement, $\overline{G}$, contains a bi-clique of size $\Omega(n)$. However, to apply one of these two theorems, either in G or in $\overline{G}$, the maximum degree of the vertices must be sufficiently small, which is not necessarily the case.

To overcome this difficulty, we use the following statement which guarantees that G or $\overline{G}$ has a large induced subgraph with very few edges.

Lemma 9. *Let $H_<$ be an ordered graph and let $\epsilon > 0$. Then there exists a constant $c_0 = c_0(H_<, \epsilon) > 0$ such that every ordered graph $G_<$ on n vertices that does not contain $H_<$ as an induced subgraph has the following property. There is a subset $U \subset V(G)$ with $|U| \geq c_0 n$ such that either $|E(G[U])| \leq \epsilon\binom{|U|}{2}$ or $|E(G[U])| \geq (1-\epsilon)\binom{|U|}{2}$ holds.*

Lemma 9 is an easy consequence of the unordered variant of the same statement due to Rödl [18] and a result of Rödl and Winkler [19].

Lemma 10. *(Theorem 1 in [18]) Let H be a graph and let $\epsilon > 0$. Then there exists a constant $c_1 = c_1(H, \epsilon) > 0$ such that every graph G on n vertices that does not contain H as an induced subgraph has the following property. There is a subset $U \subset V(G)$ with $|U| \geq c_1 n$ such that either $|E(G[U])| \leq \epsilon\binom{|U|}{2}$ or $|E(G[U])| \geq (1-\epsilon)\binom{|U|}{2}$ holds.*

Lemma 11. *[19] For every ordered graph $H_<$, there exists an unordered graph H' with the property that for any total ordering $\prec$ on $V(H')$, the ordered graph $H'_\prec$ contains $H_<$ as an induced subgraph.*

Proof of Lemma 9. Let H' be the graph whose existence is ensured by Lemma 11. For every ordered graph $G_<$ that does not contain $H_<$ as an induced subgraph, the underlying unordered graph G does not contain H' as an induced subgraph. Hence, the statement is true with $c_0 = c_1(H', \epsilon)$, where $c_1(H', \epsilon)$ is the constant defined in Lemma 10. □

Corollary 12. *Let $H_<$ be an ordered graph and let $\delta > 0$. There exists a constant $c_2 = c_2(H_<, \delta) > 0$ such that every ordered graph $G_<$ on n vertices that does not contain $H_<$ as an induced subgraph has the following property. There is a subset $U \subset V(G)$ with $|U| \geq c_2 n$ such that either $\Delta(G[U]) \leq \delta|U|$ or $\Delta(\overline{G}[U])| \leq \delta|U|$ holds.*

Proof. Let $\epsilon = \delta/4$, and let $c_0 = c_0(H_<, \epsilon)$ be the constant given by Lemma 9. We show that $c_2 = c_0/2$ meets the requirements.

Let $G_<$ be an ordered graph on n vertices that does not contain an induced copy of $H_<$. Then there exists $U_0 \subset V(G)$ with $|U_0| \geq c_0 n$ such that either $|E(G[U_0])| \leq \epsilon\binom{|U_0|}{2}$ or $|E(G[U_0])| \geq (1-\epsilon)\binom{|U_0|}{2}$ holds. Without loss of generality, suppose that $|E(G[U_0])| \leq \epsilon\binom{|U_0|}{2}$; the other case can be handled in a similar manner. Let U_1 be the set of vertices $u \in U_0$ whose degree in $G[U_0]$ is larger than $2\epsilon|U_0|$. Clearly, we have $|U_1| < |U_0|/2$. Setting $U = U_0 \setminus U_1$, we obtain $|U| > |U_0|/2 \geq c_2 n$, and the degree of every vertex in $G[U]$ is at most $2\epsilon|U_0| < 4\epsilon|U| = \delta|U|$. □

Now we are in a position to prove Lemma 7.

Proof of Lemma 7. Let $c' = c(M_1)$ be the constant defined in Theorem 2, and let $c'' = c(4)$ be the constant defined in Theorem 1. Set $\delta = \min\{c', c''\}$, and let $c_2 = c_2(M_1, \delta)$ be the constant defined in Corollary 12.

By Lemma 8 (1), there exists an ordering $<$ on G such that $G_<$ does not contain M_1 as an induced subgraph. Hence, there exists $U \subset V(G)$ such that $|U| \geq c_2 n$, and either $\Delta(G[U]) < \delta |U|$, or $\Delta(\overline{G}[U]) < \delta |U|$. In the first case, $\overline{G}[U]$ contains a bi-clique of size $c'|U| \geq c' c_2 n$. In the second case, by Lemma 8 (2), $\overline{G}_<$ does not contain P_4 as an induced subgraph, so $G[U]$ contains a bi-clique of size $c''|U|/\log|U| \geq c'' c_2 n/\log n$. Thus, the statement is true with $c = \delta c_2$. □

Next, we show that Theorem 3 holds not only for families of grounded x-monotone curves, but also for families $\mathcal{C}$ of x-monotone curves, each of which intersects the same vertical line. Clearly, such a line splits $\mathcal{C}$ into two families of grounded curves, and the intersection graph of $\mathcal{C}$ is the union of the intersection graphs of these two families. In order to exploit this property, we make use of the following technical lemma. The constants c and c' appearing in Lemma 13 are different from all previously used constants denoted by the same letters. Let us note that Theorem 8 in [7] is a result of similar flavor, but it is not suitable for our purposes.

A family of graphs $\mathcal{G}$ is called *hereditary*, if for every $G \in \mathcal{G}$, every induced subgraph of G is also a member of $\mathcal{G}$. For any pair of graphs G_1 and G_2 with $V(G_1) = V(G_2)$, the *union of* G_1 *and* G_2 is defined as the graph $G_1 \cup G_2$ whose vertex set is $V(G_1)$ and edge set is $E(G_1) \cup E(G_2)$.

Lemma 13. *Let $\mathcal{G}$ be a hereditary family of graphs. Suppose that there exist a constant $c, 0 < c < 1$, and a monotone increasing function $f : \mathbb{N} \to \mathbb{R}^+$ such that each member $G \in \mathcal{G}$ on n vertices contains either a bi-clique of size at least $n/f(n)$, or $\overline{G}$ contains a bi-clique of size at least cn.*

Then there exists a constant $c' > 0$ with the following property. If $G_1, G_2 \in \mathcal{G}$, $V(G_1) = V(G_2)$, and $|V(G_1)| = n$, then $G_1 \cup G_2$ contains a bi-clique of size at least $c'n/f(n)$ or the complement of $G_1 \cup G_2$ contains a bi-clique of size at least $c'n$.

Proof. Let $k = 1 + \lceil \log_2(1/c) \rceil$. We show that the constant $c' = c^{k+1}/2$ will meet the requirements. Let $G_1, G_2 \in \mathcal{G}$ such that $V = V(G_1) = V(G_2)$ and $|V| = n$.

We can suppose that if $U \subset V$ such that $|U| \geq \frac{c'}{c} n$, then both $\overline{G}_1[U]$ and $\overline{G}_2[U]$ contain a bi-clique of size $c|U|$. Indeed, otherwise, either $G_1[U]$ or $G_2[U]$ contains a bi-clique of size $c|U|/f(|U|) \geq c'n/f(n)$, so $G_1 \cup G_2$ also contains a bi-clique of size $c'n/f(n)$, and we are done.

For $i = 0, \dots, k$, we define disjoint sets $U_{i,1}, \dots, U_{i,2^i} \subset V$ such that $|U_{i,j}| \geq c^i n$ for $j = 1, \dots, 2^i$, and there is no edge between $U_{i,j}$ and $U_{i,j'}$ in G_1 for $1 \leq j < j' \leq 2^i$. Let $U_{0,1} = V$. If $U_{i,1}, \dots, U_{i,2^i}$ are already defined for $i < k$, let $(U_{i+1,2j-1}, U_{i+1,2j})$ be a bi-clique of size $c|U_{i,j}|$ in $\overline{G}_1[U_{i,j}]$. As $|U_{i,j}| = c^i n > \frac{c'}{c} n$, such a bi-clique always exists.

Now let $U = \bigcup_{j=1}^{2^k} U_{k,j}$. Then $|U| = 2^k c^k n > \frac{c'}{c} n$, so $\overline{G}_2[U]$ contains a bi-clique (A, B) of size at least $c|U|$. Therefore, there exists $1 \leq j \leq 2^k$ such that $|U_{k,j} \cap A| \geq |A|/2^k \geq c|U|/2^k = c^{k+1} n > c'n$, and there exists $1 \leq j' \leq k$ such that $j \neq j'$ and

$$|U_{k,j'} \cap B| \geq \frac{|B| - |U_{k,j}|}{2^k} \geq \frac{c|U| - |U|/2^k}{2^k} = \left(c^{k+1} - \frac{c^k}{2^k}\right) n \geq \frac{c^{k+1} n}{2} = c'n.$$

There is no edge between $A \cap U_{k,j}$ and $B \cap U_{k,j'}$ in $\overline{G}_1$ and $\overline{G}_2$, so the complement of $G_1 \cup G_2$ contains a bi-clique of size $c'n$. □

Now we can prove Theorem 3 for collections of x-monotone curves that intersect the same vertical line.

Lemma 14. *Let $\mathcal{C}$ be a collection of n x-monotone curves such that each member of $\mathcal{C}$ intersects a vertical line l. Let G be the intersection graph of $\mathcal{C}$. Then either G contains a bi-clique of size $\Omega(n/\log n)$, or the complement of G contains a bi-clique of size $\Omega(n)$.*

Proof. Let $\mathcal{G}$ be the family of intersection graphs of collections of grounded x-monotone curves. Clearly, $\mathcal{G}$ is hereditary. By Lemma 7, there exists a constant $c > 0$ such that each $G_0 \in \mathcal{G}$ on n vertices contains either a bi-clique of size $cn/\log n$, or the complement of G_0 contains a bi-clique of size cn. Thus, by Lemma 13, there exists a constant $c' > 0$ such that if $G_1, G_2 \in \mathcal{G}$ with $V(G_1) = V(G_2)$ and $|V(G_1)| = n$, then either $G_1 \cup G_2$ contains a bi-clique of size $c'n/\log n$, or the complement of $G_1 \cup G_2$ contains a bi-clique of size $c'n$.

The vertical line l cuts each x-monotone curve $\alpha \in \mathcal{C}$ into a left and a right part, denoted by α_1 and α_2. Let $\mathcal{C}_1 = \{\alpha_1 : \alpha \in \mathcal{C}\}$, $\mathcal{C}_2 = \{\alpha_2 : \alpha \in \mathcal{C}\}$, and let G_1 and G_2 be the intersection graphs of $\mathcal{C}_1$ and $\mathcal{C}_2$, respectively. Then $G_1, G_2 \in \mathcal{G}$ and $G = G_1 \cup G_2$, so we are done. □

Finally, everything is ready to prove our main theorem.

Proof of Theorem 3. For each $\alpha \in \mathcal{C}$, let $r(\alpha)$ denote the x-coordinate of the right endpoint of α. Without loss of generality, we can suppose that $r(\alpha) \neq r(\alpha')$ for $\alpha \neq \alpha'$. Let $\alpha_1, \ldots, \alpha_n$ be the enumeration of the curves in $\mathcal{C}$ such that $r(\alpha_1) < \cdots < r(\alpha_n)$.

Set $m = \lfloor n/3 \rfloor$ and consider a vertical line $l = \{x = r\}$, where $r(\alpha_m) < r < r(\alpha_{m+1})$. Let $\mathcal{C}'$ denote the set of curves in $\mathcal{C}$ which have a nonempty intersection with l. We distinguish two cases.

Case 1: $|\mathcal{C}'| \geq m$. Let G' be the intersection graph of $\mathcal{C}'$. Then by Lemma 14, either G' contains a bi-clique of size $\Omega(m/\log m) = \Omega(n/\log n)$, or $\overline{G}'$ contains a bi-clique of size $\Omega(m) = \Omega(n)$.

Case 2: $|\mathcal{C}'| < m$. Let $A = \{C_i : i \leq m\}$ and $B = \mathcal{C} \setminus (A \cup \mathcal{C}')$. Then $|B| \geq n/3$, and no curve in A intersects any curve in B, because A and B are separated by l. Hence, $\overline{G}$ contains a bi-clique of size $m = \Omega(n)$. □

5 Sharp threshold for intersection graphs–Proof of Theorem 4

In this section, we prove Theorem 4. Part (1) of the theorem is an easy consequence of the following result of Pach and Tóth [17]; see also [13].

Lemma 15. *(Pach, Tóth [17]) Let V be an n-element set and let V_1, V_2, V_3, V_4 be a partition of V into 4 sets. Let G be a graph on the vertex set V such that V_i spans a clique in G for $i = 1, 2, 3, 4$.*

Then G can be realized as the intersection graph of convex sets.

Proof of Theorem 4, part (1). Let V be an n-element set and let V_1, V_2, V_3, V_4 be a partition of V into four sets of size roughly $n/4$. Consider the graph G in which V_1, V_2, V_3, V_4 are cliques, and any pair of vertices $\{u, v\}$, where $u \in V_i$ and $v \in V_j$ with $i \neq j$ is joined by an edge with probability ϵ. Then with probability tending to 1, G has at most $(\frac{1}{4} + \epsilon)\binom{n}{2}$ edges, and $\overline{G}$ contains no bi-clique of size $4\frac{\log n}{\epsilon}$. By Lemma 15, G can be realized as the intersection graph of convex sets and, therefore, by x-monotone curves. □

In the rest of the section, we prove part (2) of Theorem 4. For the proof, we use the following characterization of intersection graphs of x-monotone curves that intersect the same vertical line, which was established in [15].

A graph $G_{<_1,<_2}$ with two total orderings, $<_1$ and $<_2$, on its vertex set is called *double-ordered*. If the orderings $<_1, <_2$ are clear from the context, we shall write G instead of $G_{<_1,<_2}$.

Definition 16. *A double-ordered graph $G_{<_1,<_2}$ is called* magical *if for any three distinct vertices $a, b, c \in V(G)$ with $a <_1 b <_1 c$ the following is true: if $ab, bc \in E(G)$ and $ac \notin E(G)$, then $b <_2 a$ and $b <_2 c$. A graph G is said to be* magical *if there exist two total orders $<_1, <_2$ on $V(G)$ such that $G_{<_1,<_2}$ is magical.*

A *triple-ordered* graph is a graph $G_{<_1,<_2,<_3}$ with three total orders $<_1, <_2, <_3$ on its vertex set.

Definition 17. *A* triple-ordered graph $G_{<_1,<_2,<_3}$ *is called* double-magical*, if there exist two magical graphs $G^1_{<_1,<_2}$ and $G^2_{<_1,<_3}$ on $V(G)$ such that $E(G_{<_1,<_2,<_3}) = E(G^1_{<_1,<_2}) \cap E(G^2_{<_1,<_3})$. An* unordered graph G *is said to be* double-magical *if there exist three total orders $<_1, <_2, <_3$ on $V(G)$ such that the triple-ordered graph $G_{<_1,<_2,<_3}$ is double-magical.*

Lemma 18. *(Pach, Tomon [15]) A graph is double-magical if and only if it is isomorphic to the complement of the intersection graph of a collection of x-monotone curves, each of which intersects a vertical line l.*

Here, we will only use the easier direction that if G is the complement of the intersection graph of a collection $\mathcal{C}$ of x-monotone curves, each of which intersects a vertical line l, then G is double-magical. To see why this is true, let $<_1$ be the ordering given by the intersection points of the elements of $\mathcal{C}$ with the vertical line l, and let $<_2$ and $<_3$ be the orderings induced by the x-coordinates of the right and left endpoints of the curves, respectively. Then it is not hard to show that $G_{<_1,<_2,<_3}$ is double-magical.

Relying on the characterization provided by Lemma 18, part (2) of Theorem 4 reduces to the following lemma about double-magical graphs.

Lemma 19. *For any $\epsilon > 0$, there exists a constant $c = c(\epsilon) > 0$ with the following property. For every positive integer n, every double-magical graph with n vertices and at least $(\frac{3}{4} + \epsilon)\binom{n}{2}$ edges contains a bi-clique of size cn.*

Proof. Let $G_{<_1,<_2}$ and $G_{<_1,<_3}$ be magical graphs such that $E(G) = E(G^1_{<_1,<_2}) \cap E(G^2_{<_1,<_3})$.

A triple of vertices (a, b, c) in G is called an *i-hole* for $i = 2, 3$, if $a <_1 b <_1 c$ and $b <_i a$ and $b <_i c$. A 4-tuple (a, b, b', c) of vertices of G is said to be *forcing* if

1. $a <_1 b <_1 c$,
2. $a <_1 b' <_1 c$,
3. (a, b, c) is not a 2-hole,
4. (a, b', c) is not a 3-hole.

Note that we do not exclude that $b = b'$. If (a, b, b', c) is forcing, we say that the *set* $\{a, b, b', c\}$ is also *forcing*. We are interested in forcing 4-tuples for the following reason: if (a, b, b', c) is forcing such that $ab, bc, ab', b'c$ are edges of G, then ac is also an edge. Indeed, if ab, bc are edges of G, then ab, bc are edges of $G_{<_1,<_2}$. In this case, as $G_{<_1,<_2}$ is magical and (a, b, c) is not a 2-hole, ac is also an edge of $G_{<_1,<_2}$. Similarly, $ab', b'c$ are edges of $G_{<_1,<_3}$. Then, as $G_{<_1,<_3}$ is magical and (a, b', c) is not a 3-hole, ac is also an edge of $G_{<_1,<_3}$. Therefore, ac is an edge of G as well.

Let the *order type* of a 4-tuple (a, b, b', c) of vertices be (s_1, s_2, s_3, s_4), where

$$s_1 = \begin{cases} + & \text{if } a <_2 b \\ - & \text{if } b <_2 a \end{cases}, \quad s_2 = \begin{cases} + & \text{if } b <_2 c \\ - & \text{if } c <_2 b \end{cases}, \quad s_3 = \begin{cases} + & \text{if } a <_3 b' \\ - & \text{if } b' <_3 a \end{cases}, \quad \text{and} \quad s_4 = \begin{cases} + & \text{if } b' <_3 c \\ - & \text{if } c <_3 b' \end{cases}.$$

Note that if (a, b, b', c) is a 4-tuple such that $a <_1 b <_1 c$ and $a <_1 b' <_1 c$, and the order type of (a, b, b', c) is (s_1, s_2, s_3, s_4), then (a, b, b', c) is not forcing if and only if $(s_1, s_2) = (-, +)$, or $(s_3, s_4) = (-, +)$.

Claim 20. *Every set of* 5 *vertices in* G *contains a forcing* 4*-tuple.*

Proof. There are $(5!)^2 = 14400$ non-isomorphic triple orderings of a 5 elements set, so it is sufficient to show that each of them contains a forcing 4-tuple. A quick computer search shows that this is indeed the case. We provide a more detailed proof in the Appendix, which reduces the number of cases to a 120. □

As usual, let K_t denote the complete graph on t vertices. By a well known result of Erdős and Simonovits [5], the condition $|E(G)| \geq (1 - \frac{1}{4} + \epsilon)\binom{n}{2}$ implies that G contains at least $c_0 n^5$ copies of the complete graph K_5, where $c_0 = c_0(\epsilon)$ depends only on ϵ.

Now each copy of K_5 in G contains a forcing 4-tuple by Claim 20, which spans either a copy of K_4 or K_3 in G (depending on whether $b = b'$). There are 16 order types of 4-tuples. Hence, there is an order type τ such that either at least $c_0 n^5/32$ copies of K_5 in G contain a copy of K_4 that is forcing with order type τ, or at least $c_0 n^5/32$ copies of K_5 in G contain a copy of K_3 that is forcing with order type τ. As every copy of K_4 is contained in at most n copies of K_5, and every copy of K_3 is contained in at most n^2 copies of K_5, we get the following two cases. Either there exist at least $c_0 n^4/32$ copies of K_4 that is forcing with order type τ, or there exist at least $c_0 n^3/32$ copies of K_3 in G that is forcing with order type τ.

In the first case, we deduce that there exists a pair of vertices (b, b') in G such that $b \neq b'$, and there are at least $c_0 n^2/32$ pairs of vertices (a, c) such that (a, b, b', c) is forcing with order type τ, and $\{a, b, b', c\}$ spans a copy of K_4. Let A be the set of vertices a that appear in such a forcing 4-tuple (a, b, b', c), and let C be the set of vertices c that appear in such a forcing 4-tuple (a, b, b', c). Then $|A||C| \geq c_0 n^2/32$, so $|A|, |C| \geq c_0 n/32$. If $a_0 \in A$, there exists $c \in C$ such that $\{a_0, b, b', c\}$ spans a copy of K_4, so a_0 is joined to b and b' by an edge. Similarly, every $c_0 \in C$ is also connected to b and b' by an edge. Finally, for every $a_0 \in A$ and $c_0 \in C$, the 4-tuple (a_0, b, b', c_0) has order type τ. But whether a 4-tuple is forcing depends only on its order type, so (a_0, b, b', c_0) is forcing. But then $a_0 c_0$ is an edge, otherwise, either $b <_2 a_0$ and $b <_2 c_0$, or $b' <_3 a_0$ and $b' <_3 c_0$ holds by Definition 16, which implies that (a_0, b, b', c_0) is not forcing, contradiction. In conclusion, a_0c_0 is an edge for every $a_0 \in A$ and $c_0 \in C$, so $A \cup C$ spans a bi-clique of size at least $c_0n/32$.

In the second case, there exist a vertex b and at least $c_0n^2/32$ pairs of vertices (a, c) such that (a, b, b, c) is forcing with order type τ, and $\{a, b, c\}$ spans a copy of K_3. Now we can proceed in the same way as in the previous case to find a bi-clique of size $c_0n/32$.

Hence, Lemma 19 holds with $c = c_0/32$. □

Corollary 21. *For any $\epsilon > 0$, there exist a constant $c = c(\epsilon) > 0$ and an integer $n_0 = n_0(\epsilon)$ such that the following statement is true. For any $n \geq n_0$ x-monotone curves that intersect the same vertical line, if the intersection graph G of the curves has at most $(\frac{1}{4} - \epsilon)\binom{n}{2}$ edges, then the complement of G contains a bi-clique of size cn.*

Proof. By Lemma 18, the complement of G is a double-magical graph. Since $\overline{G}$ has at least $(\frac{3}{4}+\epsilon)\binom{n}{2}$ edges, by Lemma 19 it must contain a bi-clique of size cn. □

Similarly as in the proof of Theorem 3, we complete the proof of Theorem 4 by reducing the general configuration of x-monotone curves to the case, where every x-monotone curve has nonempty intersection with the same vertical line l.

Proof of Theorem 4, part (2). Without loss of generality, assume that $\epsilon < 1/2$. Let $\mathcal{C}$ denote our collection of curves. For each $\alpha \in \mathcal{C}$, let $r(\alpha)$ be the x-coordinate of the right endpoint of α. We can also suppose that $r(\alpha) \neq r(\alpha')$ for $\alpha \neq \alpha'$. Let $\alpha_1, \ldots, \alpha_n$ be the enumeration of the curves in $\mathcal{C}$ such that $r(\alpha_1) < \cdots < r(\alpha_n)$.

Set $m = \epsilon n/2$ and consider a vertical line $l = \{x = r\}$, where $r(\alpha_m) < r < r(\alpha_{m+1})$. Let $\mathcal{C}'$ denote the set of curves in $\mathcal{C}$ which have a nonempty intersection with l. We distinguish two cases.

Case 1: $|\mathcal{C}'| \geq (1-\epsilon)n$. Let G' be the intersection graph of $\mathcal{C}'$. Then G' has at most $\left(\frac{1}{4} - \frac{\epsilon}{4}\right)\binom{|V(G')|}{2}$ edges. Therefore, the complement of G' contains a bi-clique of size $c(\frac{\epsilon}{4})|V(G')| > c(\frac{\epsilon}{4})n/2$, where c is the constant defined in Corollary 21.

Case 2: $|\mathcal{C}'| < (1-\epsilon)n$. Let $A = \{C_i : i \leq m\}$ and $B = \mathcal{C} \setminus (A \cup \mathcal{C}')$. Then $|B| \geq \epsilon n/2$, and no curve in A intersects any curve in B, because A and B are separated by l. Hence, $\overline{G}$ contains a bi-clique of size $\epsilon n/2$. □

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Appendix - Proof of Claim 20

Let $<_1$ denote the natural total ordering on $\{1,\dots,5\}$. We need to prove that for any two total orderings $<_2$ and $<_3$ on $\{1,\dots,5\}$, there exists a 4-tuple (a,b,b',c) in $\{1,\dots,5\}$ such that $a <_1 b <_1 c$, $a <_1 b' <_1 c$, (a,b,c) is not a 2-hole, and (a,b',c) is not a 3-hole.

If $\prec$ is a total ordering on $\{1,\dots,5\}$, let $H(\prec)$ denote the graph on vertex set $\{1,\dots,5\}$, where $a <_1 c$ are joined by an edge if there exists b such that $a <_1 b <_1 c$, and b satisfies at least one of the inequalities $a \prec b$ and $c \prec b$.

Our task is reduced to proving $E(H(<_2)) \cap E(H(<_3)) \neq \emptyset$. Indeed, if $ac \in E(H(<_2)) \cap E(H(<_3))$, then there exists b, b' such that (a,b,b',c) is a desired 4-tuple. Say that a graph H is *attainable*, if there exists a total ordering $\prec$ such that $H(\prec) = H$. We want to show that if H and H' are both attainable, then $E(H) \cap E(H') \neq \emptyset$. Note that if a graph is attainable, then its edge set is a subset of the 6 element set $\{13, 14, 15, 24, 25, 35\}$. Using a quick computer search on all the possible 120 total orderings $\prec$ (which might also have been done by hand, if one is patient enough), we determined all the possible attainable graphs. We found that each of the attainable graphs has at least 3 edges, so if $E(H)$ and $E(H')$ does not intersect, then we must have $|E(H)| = |E(H')| = 3$. In total, there are 5 attainable graphs with 3 edges:

$$\{13, 14, 24\}, \{13, 14, 35\}, \{14, 24, 25\}, \{13, 25, 35\}, \{24, 25, 35\}.$$

It is easy to check that any two of these edge sets have a nonempty intersection.