G22.1170: Fundamental Algorithms<br>Final Take-Home Exam<br>(Due Tuesday December 19 2000)

Problem. 1 Order the following functions in increasing order of growth:
(a) $n$ !
(b) $\lg \lg n$
(c) $n^{\lg \lg n}$
(d) $n^{1 / \lg n}$
(e) $2^{3 \lg n}$
(f) $n^{2}$
(g) $(\lg \lg n)^{n}$
(h) $n \lg ^{2} n$

Problem. 2 The input is a sequence of $n$ elements $x_{1}, x_{2}, \ldots, x_{n}$ that we can read sequentially. We want to use a memory that can only store $O(k)$ elements at a time. Give a high level description of an algorithm that finds the $k$ th smallest element in $O(n)$ time.

Problem. 3 Let $L$ be a sequence of $n$ elements. If $x$ and $y$ are pointers into list $L$ then $\operatorname{Insert}(x)$ inserts a new element immediately to the right of $x$, $\operatorname{Delete}(x)$ deletes the element to which $x$ points and $\operatorname{Order}(x, y)$ returns true if $x$ is before $y$ in the list. Show how to implement all three operations with worst case time $O(\log n)$.

Problem. 4 A simple undirected graph $G=(V, E)$ without self-loops has at most one edge between every pair of vertices and no edge from a vertex to itself. A graph is $p$-colorable if all vertices can be assigned one of $p$ colors with no edge receiving the same color at both of its ends.

Let $d(v)$ denote the degrees of a vertex $v$, i.e., the number of edges incident at $v$. let $d(G)$ denote $\max _{v \in V} d(v)$, the maximum degree of the vertices of the graph $G$.

Design an efficient algorithm and prove its correctness, which determines $(d(G)+1)$-coloring of the graph.

