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Lecture # 9

—Slide 1—

The ADA Programming Language Language Survey 3

- Countess Ada Augusta Lovelace
 First programmer (hacker) on Babbage's Analytical
 Engine.
- Relatively new programming Language (1980) developed by DoD. Currently, the only language approved for DoD software project
- Developed to reduce high cost of designing, writing and maintaining DoD software.
- Major items of software are for "**embedded systems**"
 - Primary purpose is control.
 - Incorporated in larger systems.
 - Large software: 50,000-100,000 lines of code
 - Long-lived systems (10–15 years)
 - Continuously evolving.

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The Design History

- 1975: HOLWG (High Order Language Working Group) was established to investigate problems of developing common language for DoD.
- 1975–1978: List of requirements: STRAWMAN, WOODENMAN, TINMAN, IRONMAN, Revised IRONMAN, and finally STEELMAN
- Design Competition: Winner: Jean Ichbiah, Cii Honeywell Bull
- 1987: Ada ISO Standard 8652
- 1983–1987: Analysis of Technical Queries by ARG (Ada Rapporteur Group). Resulted in AIs (Ada Issues).
- 1988: DOD (and ANSI) established Ada 9X project.
- 15 Feb 1995: Ada 95. Core Language + Specialized Annexes

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Quick Overview (Ada 83)

- Strong Typing
 Scalar, Composite, Access, Private & Derived
- Representation Specification
- Standard Control Constructs: Structured Language—Pascalish
- Subprograms: Procedures & Functions
- Program Structuring Facilities: Packages
- Generics: Polymorphisms Extends the concept of type and functional abstractions
- Exceptions
- Separate Compilation
- Tasking (Rendezvous)

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Quick Overview (Ada 95)

- Object Oriented Programming
- Program Libraries
 (Replaces Ada 83's flat structure by a hierarchy—Visibility
 Control and Program Extension without recompilation)
- Interfacing
 Graphical User Interface (GUI)—Program Call-back.
- Tasking (Static Monitor)

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Ada Type System

- Programmer Defined Types
 Set of values, Applicable set of operations and set of properties, accessible via attributes
- Type Definition

type <identifier> is <type-definition>

• Type Binding

<variable> : <type>

• Attributes

<name>'<attribute-identifier>

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Examples

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Primitive Scalar Types

- Discrete Types
 - i) Enumeration Types, ii) Character type,
 - iii) Boolean Types, iv) Integer Types
- Real Types
 - i) Fixed-Point Types, ii) Floating-Point Types
- **Note:** Integer and Real types together form the numeric types.
- INTEGER = Predefined integer types SYSTEM.MIN_INT..SYSTEM.MAX_INT
- Operation

+, -, *, / (Arithmetic operations)

<, <=, =, /=, >=, > (Logical operations)

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Discrete Types

• Note: Characters & Boolean are predefined enumeration types

Characters = 128 ASCII characters: 'A', 'B', 'C', ...

Booleans = FALSE & TRUE

- All discrete types are **ordered**: FALSE < TRUE.
- Discrete type values can be used for **indexing** & **iteration over loops**.
- Attributes of Discrete Types (e.g., T)

T'POS: Position Number

SUIT'POS(HEARTS) = 0

T'VAL: Inverse of POS

SUIT'VAL(0) = HEARTS

T'SUCC, T'PRED

T'FIRST, T'LAST

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Real Type

• Floating Point:
Accuracy is fixed by a *relative* error bound

type WEIGHT is digits 10;

Values have an accuracy of 10 digits.

• Fixed Point:

Accuracy is fixed by a relative error bound

type VOLTAGE is delta 0.1 range 0.0..1.0;

Values have an accuracy at least as fine as 0.1.

• Some Attributes

$$\mathbf{FX} = \text{Fixed Point: } (-2^N + 1)\delta..(2^N + 1)\delta$$

FX'DELTA, FX'LARGE

 $\mathtt{FL} = \mathtt{Floating} \ \mathtt{Point:} \ \mathtt{sgn} * \mathtt{mantissa} * 2^E; \ \mathtt{mantissa} \ \mathtt{has} \ B \ \mathtt{digits.}$

FL'DIGITS, FL'MANTISSA, FL'EMAX, FL'SMALL, FL'LARGE, FL'EPSILON

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Derived Type

- Specific Type + Constraints E.g., Range Constraints for scalar
- Examples

```
subtype BYTE_SIZE is INTEGER range -2**7..2**7;
subtype CAPS is CHARACTER range 'A'..'Z';
subtype <identifier> is <parent-type> range <constraint>;
```

• Object Declaration

• Range constraint can be tested at runtime

```
if CURRENT_INPUT not in CAPS then ...
```

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Assignment Statement

X := Y;

• Types must match

- Same type name-Not structure. Ada uses *name* equivalence.
- Type checking is at compile time.
- Type mismatch \Rightarrow Program considered *illegal*.

• Subtype constraints

- Compatibility condition: $type(X) \ge type(Y)$
- Constraint checking is at run time.
- Violation raises constraint exception.(Program is **not** illegal.)

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Examples

```
subtype NATURAL is INTEGER range 1..INTEGER'LAST;
A: INTEGER;
B: FLOAT;
C: NATURAL;
D: INTEGER range 0..INTEGER'LAST;
A := B;
                    --illegal
A := INTEGER(B); --type conversion, legal
A := C;
                    --legal
A := D - 3;
                    --legal
A := C + INTEGER(B); --legal
C := D;
                     --constraint exception
C := A;
                     --constraint exception
```

[End of Lecture #9]