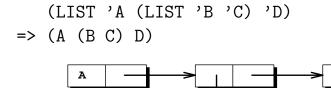
V22.0490.001 Special Topics: Programming Languages

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Lecture # 12

-Slide 1-LISTS

• Function LIST



(A (B C) D)

NIL

NIL

(LIST 'B NIL) => (B NIL)



(LIST)

=> NIL

(An Atom not a CONS cell)

—Slide 2—

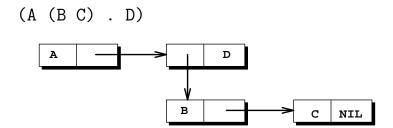
Proper and Improper Lists

• Proper Lists

Lists terminating in NIL

• Improper Lists

Lists not terminating in NIL



—Slide 3—

List Operations

- CAR: Extracts the *first* element of a list
- CDR: Extracts the *rest* (all but the first element) of a list
- CAR & CDR can be applied to any list, but not to atom other than NIL

$$(CAR\ NIL) \equiv (CDR\ NIL) \equiv NIL$$

• Examples

—Slide 4—

List Predicates

• Boolean-valued functions

$${T = True, NIL = False}$$

• ATOM: True iff an atom

$$(ATOM 'NIL) => T$$

 $(ATOM '(X Y)) => NIL$

- CONSP, LISTP, ...
- NULL: True iff an empty list (e.g., NIL)

```
(NULL 'NIL) => T
(NULL 'X) => NIL
```

• ZEROP, NUMBERP, ...

—Slide 5—

Conditional Form

• IF & COND

• IF form is equivalent to

- When the LISP evaluation procedure sees a COND, it evaluates each of the test forms in order (i.e., <test-1>,</test-2>,
- When one evaluates to "non-NIL," its corresponding result form (i.e., <result-i>) is evaluated and its value is returned as the result of the COND.

—Slide 6—

Simple Examples

• Example

```
(IF (ATOM 'X) 'YES 'NO)
                                     => YES
(COND ((ATOM 'X) 'YES)
      (T 'NO))
                                        YES
(DEFUN LENGTH (LIST)
      (IF (NULL LIST)
          0
          (+ 1 (LENGTH (CDR LIST)))))
```

• LENGTH Computation

—Slide 7—

LISP EVALUATOR: EVAL

LISP Obj \Rightarrow S-expression (Symbolic Expn) \mapsto Values

Evaluation Order = "Inside-Out"

Atoms evaluate to themselves:
 Numbers and NIL are constants and have predefined values. Symbols can be SET to a value in the course of execution.

$$(SET 'A '(X Y)) => A = (X Y)$$

• List to be evaluated must have a "functional-valued" symbol as its CAR. The CDR of the list is treated as the arguments.

Evaluation proceeds by applying the function (CAR) to the argument (CDR)

$$(+ 2 4)$$
 => 6
 $(+ (CAR (CDR (LIST 2 3))) 4)$ => 7

—Slide 8—

QUOTE

• Evaluation Blocker

- \circ If the symbol in the functional position is a *special form*, then the usual evaluation rule is replaced by a special one *before* the arguments are evaluated
- QUOTE \equiv Special Form

A quoted item evaluates to itself

• Example

—Slide 9—

EVAL

• Evaluation procedure can be invoked directly

```
(SETQ A (LIST '+ 5 6)) => A = (+ 5 6)

(CONS A '(IS MY RESULT))

=> ((+5 6) IS MY RESULT)

(CONS (EVAL A) '(IS MY RESULT))

=> (11 IS MY RESULT)
```

—Slide 10— USER-DEFINED FUNCTIONS

(DEFUN <function-name> <argument-list> <body>)

• DEFUN

Special Form: subforms are not evaluated as arguments

• <function-name>

Must be a symbol. Can be any symbol other than NIL

• <argument-list>

Number of Symbols in the argument list is the number of arguments the function must be handed by the **EVAL** procedure

• <body>

An S-expression that will be evaluated when the function is called... After appropriate substitution of the actuals for formals.

—Last Slide—

$Examples: APPEND \ \mathcal{E} \ REVERSE$

```
(DEFUN APPEND (X Y)

(IF (NULL X)

Y

(CONS (CAR X) (APPEND (CDR X) Y))))

(DEFUN REVERSE (LIST)

(IF (NULL LIST)

'NIL

(APPEND (REVERSE (CDR LIST))

(CONS (CAR LIST) 'NIL))))
```

[End of Lecture #12]