Sipser text, no. 5.1 (2nd edition), 5.1 (1st edition). A string $\langle G_1, G_2 \rangle$ in this language represents two CFGs $G_1$ and $G_2$ that generate the same language.  

Hint: Assume for a contradiction that there is a decision TM $R$ for $EQ_{CFG}$, and use it to build a decision TM $S$ for $ALL_{CFG}$, which is a contradiction, as we showed in class that there is no such TM $S$.

2. Show that $ALL_{TM} = \{ \langle M \rangle \mid M \text{ is a TM that accepts all strings in its input alphabet} \}$ is undecidable.

Hint: Use a construction similar to Theorem 5.2 which shows that $E_{TM}$ is undecidable.

3. Show that $ZERO_{TM} = \{ \langle M \rangle \mid M \text{ is a TM with input alphabet } \{0, 1\} \text{ that accepts input 0} \}$ is undecidable.

Hint: Again, given a decision TM $R$ for $ZERO_{TM}$, build a TM $S$ that decides $A_{TM}$. The approach of Theorem 5.2 is helpful. That is, $S$ first builds (computes a description of) a TM $\tilde{M}_y$ that does something appropriate. Then $S$ simulates $R$ on input $\tilde{M}_y$. 
